# How many smooth numbers and smooth polynomials are there? ViBrANT Seminar, May 2, 2023

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#### 1 Definition and motivation

A positive integer n is said to be y-smooth if its primes factors do not exceed y:  $p \mid n \implies p \leq y$ . The talk will be concerned with the counting function

$$\Psi(x, y) := \#\{n \le x : n \text{ is } y \text{-smooth}\}.$$

Note  $\Psi(x,x) = \lfloor x \rfloor$ ,  $\Psi(x,1) = 1$  and  $\Psi(x,2) = 1 + \lfloor \log_2 x \rfloor$ , and that the indicator function of y-smooth numbers is completely multiplicative.

One can define an analogous quantity in the polynomial setting. A polynomial  $f \in \mathbb{F}_q[T]$  is said to be *m*-smooth if its irreducible factors have degrees bounded by  $m: P \mid f \implies \deg(P) \leq m$ . The talk will be focused today mostly on  $\Psi(x, y)$ .

Smooth numbers play an important role in cryptography. Pomerance, in the 80s, devised his Quadratic Sieve, an algorithm that (heuristically) factors integers in subexponential time, namely n is factored in  $\exp((\log n)^{1/2+o(1)})$  time. We describe it (in a loose way) below.

For i = 1, 2, ... we do the following. We take  $x_i := \lfloor \sqrt{n} \rfloor + i$ , square it and reduce it modulo n to obtain a number  $y_i$  in [0, n-1]:

$$x_i^2 \equiv y_i \mod n.$$

We then check whether  $y_i$  is T-smooth – this can be done in O(T) operations obviously, but happens quite rarely: with probability  $(\Psi(N,T)/N)^{-1}$  the number  $y_i$  will be T-smooth (heuristically). When it is T-smooth, we obtain a relation of the form

$$x_i^2 \equiv \prod_{p \le T} p^{e_{i,p}} \bmod n.$$

We want to obtain T such relations, which takes  $T^2 \times (\Psi(N,T)/N)^{-1}$  operations. Then we can perform Gaussian elimination on the T binary vectors  $\{(e_{i,p2} \mod 2)_{p \leq T}\}_{i \in S}$  where S corresponds to  $y_i$  that are T-smooth. The complexity of Gaussian elimination is  $T^3$ . It finds subset(s)  $S' \subseteq S$  such that

$$\sum_{i \in S'} (e_{i,p})_{p \le T} \equiv 0 \mod 2$$

as vectors in  $\prod_{p < T} \mathbb{F}_2$ . This means

$$\prod_{i \in S'} x_i^2 \equiv \prod_{p \le T} p^{2b_p} \bmod n$$

for  $b_p = \sum_{i \in S'} e_{i,p}/2$ . Given a relation  $A^2 \equiv B^2 \mod n$  we can compute gcd(A - B, n) and hope to find one the factors of n.

The complexity of this algorithm is  $T^2 \times (\Psi(N,T)/N)^{-1} + T^3$ , and is minimized when

$$T \approx N/\Psi(N,T)$$

which turn out to be solved for

$$T = \exp((\log N)^{1/2 + o(1)})$$

which is also the total complexity.

This uses the relation  $\Psi(N,T) \sim N\rho(\log N/\log T)$  which was established in a wide range by Hildebrand, where  $\rho$  is the Dickman function, which we discuss next.

#### 2 The Dickman function

The function  $\rho: [0, \infty) \to (0, \infty)$  was introduced by Dickman. It has initial conditions  $\rho(u) = 1$  for  $u \in [0, 1]$ . For larger u it is defined via delay-differential equation:

$$u\rho'(u) + \rho(u-1) = 0$$
, or  

$$\rho(u) = u^{-1} \int_0^1 \rho(u-t) dt.$$

It is decreasing, and in fact we see it decreases rapidly:

$$\rho(u) \le u^{-1}\rho(u-1) \implies \rho(u) \le \Gamma(u+1)^{-1} = u^{-u(1+o(1))}.$$

Dickman proved (30s) that  $\Psi(x, y) \sim x\rho(\log x/\log y)$  for  $x \ge y \ge x^{\varepsilon}$ .

De Bruijn (50s) worked out precise asymptotics for  $\rho(u)$ . To explain them we need to introduce the Laplace transform of  $\rho$ :

$$\hat{\rho}(s) := \int_0^\infty e^{-st} \rho(t) dt.$$

De Bruijn showed

$$\hat{\rho}(s) = \exp\left(\gamma + \int_0^{-s} \frac{e^t - 1}{t} dt\right).$$

A short proof of this follows from differentiating  $\hat{\rho}(s)$  under the integral sign:

$$\hat{\rho}'(s) = -\int_0^\infty t e^{-st} \rho(t) dt = -\int_0^1 t e^{-st} dt - \int_1^\infty (\int_{t-1}^t \rho(v) dv) e^{-st} dt$$
$$= -\int_0^\infty \rho(v) (\int_v^{v+1} e^{-st} dt) dv = \frac{e^{-s} - 1}{s} \hat{\rho}(s).$$

(This determines  $\hat{\rho}$  up to a multiplicative constant; see de Bruijn's work for working out the constant.) For any  $c \in \mathbb{R}$  we have

$$\rho(u) = \frac{1}{2\pi i} \int_{(-c)} e^{su} \hat{\rho}(s) ds.$$

We choose c so that  $e^{-cu}\hat{\rho}(-c)$  is minimized, i.e. c is the minimizer of

$$c \mapsto -cu + \gamma + \int_0^c \frac{e^t - 1}{t} dt.$$

Differentiating (with respect to c) we find

$$-u + \frac{e^c - 1}{c} = 0$$

So the optimal c is  $\xi(u)$  (a function of u) where  $\xi(u) \sim \log u$  is defined implicitly via

$$\frac{e^{\xi} - 1}{\xi} = u$$

Let us write

$$\rho(u) = \frac{1}{2\pi i} \int_{(-\xi(u))} e^{su} \hat{\rho}(s) ds = e^{-\xi(u)u} \hat{\rho}(-\xi(u)) \frac{1}{2\pi} \int_{\mathbb{R}} G(t) dt$$

for

$$G(t) = e^{itu}\hat{\rho}(-\xi(u) + it)/\hat{\rho}(-\xi(u)).$$

By construction G(0) = 1. By definition of  $\xi$ , G'(0) = 0. It is not hard to approximate G(t) as  $e^{-ut^2(1+o(1))/2}$  for small t (details omitted; u(1+o(1)) arises from  $(\log G)''(0)$ ). We expect

$$\rho(u) = \frac{1}{2\pi i} \int_{(-\xi(u))} e^{su} \hat{\rho}(s) ds \sim e^{-\xi(u)u} \hat{\rho}(-\xi(u)) \frac{1}{2\pi} \int_{\mathbb{R}} e^{-ut^2/2} dt \sim \frac{e^{-\xi(u)u} \hat{\rho}(-\xi(u))}{\sqrt{2\pi u}}$$

and this asymptotic relation was established rigorously by de Bruijn. The quantity  $-\xi(u)$  is called the *saddle* point for  $\rho(u)$ .

### 3 Hildebrand's work

Let

$$u = \frac{\log x}{\log y}.$$

Hildebrand (80s) proved the following:

$$\Psi(x,y) = x\rho(u)\left(1 + O\left(\frac{\log(u+1)}{\log y}\right)\right)$$

holds for  $x \ge y \ge \exp((\log \log x)^{5/3+\varepsilon})$ . Under RH he showed that

$$\Psi(x,y) = x\rho(u) \exp\left(O\left(\frac{\log(u+1)}{\log y}\right)\right)$$
(3.1)

holds for  $y \ge (\log x)^{2+\varepsilon}$ . Note this does not give an asymptotic formula for  $y = (\log x)^C$ .

These two results admit alternative proofs due to Saias (80s). Hildebrand used a physical space argument while Saias used Dirichlet series and complex analysis.

Two questions that were asked:

- 1. (Hildebrand) Can one show the asymptotic relation (3.1) fails for  $y \leq (\log x)^{2-\varepsilon}$ ?
- 2. (Pomerance) Is it true that  $\Psi(x, y) \ge x\rho(u)$  for all  $x/2 \ge y \ge 2$ ? (Intuition: for very large y, there is a lower order term in  $\Psi(x, y) x\rho(u)$ , found by de Bruijn, which is positive. Moreover,  $x\rho(u) \le \Psi(x, y)$  for  $y \le \log x$  trivially since  $\Psi(x, y) \ge 1$ ,  $x\rho(u) < 1$ .)

**Theorem 3.1** (G., 2022). Fix  $\varepsilon > 0$ . Unconditionally, there are sequences  $x_n, y_n \to \infty$  such that

$$y_n = (\log x_n)^{2-\varepsilon+o(1)}$$

and

$$\frac{\Psi(x_n, y_n)}{x_n \rho(\log x_n / \log y_n)} = \exp((\log x_n)^{\varepsilon + o(1)}).$$

**Theorem 3.2** (G., 2022). Under RH, for  $(\log x)^{1+\varepsilon} \leq y \leq (\log x)^{2-\varepsilon}$  we have

$$\frac{\Psi(x,y)}{x\rho(\log x/\log y)} = \exp\left(\Theta\left(\frac{(\log x)^2}{y\log y}\right)\right).$$

An analogue of Theorem 3.2 holds unconditionally for polynomials. Here  $\Theta(f)$  stands for a function g such that  $C \ge g/f \ge c > 0$  for some positive constants C, c.

**Theorem 3.3** (G., 2022). 1. Unconditionally,  $\Psi(x, y) \ge x\rho(u)$  holds outside of

$$y \in [\log x \exp((\log \log x)^{3/5-\varepsilon}), \exp((\log \log x)^{5/3+\varepsilon})].$$

2. Under RH,  $\Psi(x,y) \ge x\rho(u)$  holds outside of

$$y \in [(\log x)^{2-\varepsilon}, (\log x)^{2+\varepsilon}]$$

3. Assume RH. If  $\psi(y) := \sum_{n \leq y} \Lambda(n) \sim y$  satisfies  $\psi(y) - y = o(\sqrt{y} \log y)$  then  $\Psi(x, y) \geq x\rho(u)$  holds for  $y \in [(\log x)^{2-\varepsilon}, (\log x)^{2+\varepsilon}]$ . Some intuition comes from the relation

$$\Psi(x,y) \sim x\rho(u)(-\zeta(1/2)\sqrt{2})\exp\left(\frac{\psi(y)-y}{\sqrt{y}\log y}\right)$$

for  $y = (1 + (\log x)/2)^2$  (RH is used in the derivation of this relation as well).

4. If RH fails, and  $\Theta > 1/2$  is the supremum of the real parts of zeros of  $\zeta$ , then for any  $\beta \in (1 - \Theta, \Theta)$ there are sequences  $x_n, y_n$  with  $y_n = (\log x_n)^{1/(1-\beta)+o(1)}$  such that

$$\Psi(x_n, y_n) < x_n \rho(\log x_n / \log y_n) \exp(-y_n^{\Theta - \beta - \varepsilon}).$$

#### 4 First oscillation result

The rest of the talk will concentrate on Theorem 3.1 and the last part of Theorem 3.3.

Let us start with the last part of Theorem 3.3.<sup>1</sup> Rankin (30s) observed that

$$\Psi(x,y) \le x^c \zeta(c,y)$$

for any c > 0, where  $\zeta(c, y) = \prod_{p \le y} (1 - p^{-c})^{-1}$  is the partial zeta function. The optimal c, that minimizes the RHS, is denoted  $\alpha = \alpha(x, y)$ :

$$\Psi(x,y) \le x^{\alpha}\zeta(\alpha,y) = \min_{c>0} x^{c}\zeta(c,y).$$

Recall also that

$$\rho(u) \sim \frac{e^{-\xi(u)u}\hat{\rho}(-\xi(u))}{\sqrt{2\pi u}}$$

Our aim is to 'marry' two classical ideas: saddle point analysis and Landau's Oscillation result (the same result that allows one to deduce  $\psi(y) - y = \Omega_{\pm}(y^{\Theta - \varepsilon})$ ).

We introduce

$$\beta = \beta(x, y) := 1 - \xi(u) / \log y$$

where  $u = \log x / \log y$ , which allows us to rewrite

$$x\rho(u) \sim \frac{x^{\beta}\hat{\rho}(\log y(\beta-1))}{\sqrt{2\pi u}}.$$

Now let's divide  $\Psi(x, y)$  by  $x\rho(u)$ :

$$\frac{\Psi(x,y)}{x\rho(u)} \ll \sqrt{u} \frac{x^{\alpha}\zeta(\alpha,y)}{x^{\beta}\hat{\rho}(\log y(\beta-1))}$$

Here is a trivial (but new) observation. Since  $\alpha$  minimizes the numerator we trivially have

$$\frac{\Psi(x,y)}{x\rho(u)} \ll \sqrt{u} \frac{x^{\beta} \zeta(\beta,y)}{x^{\beta} \hat{\rho}(-\xi(u))} = \sqrt{u} \frac{\zeta(\beta,y)}{\hat{\rho}(-\xi(u))}$$

Letting

$$F(s,y) := \log \zeta(s,y) - \log \hat{\rho}(\log y(s-1)),$$

we see

$$\frac{\Psi(x,y)}{x\rho(u)} \ll \sqrt{u}e^{F(\beta,y)}.$$

By an earlier computation,

$$\log \hat{\rho}(\log y(s-1)) = \gamma + I((1-s)\log y).$$

As for  $\log \zeta(s, y)$ , we find

$$\log \zeta(s,y) = \sum_{p \le y} -\log(1-p^{-s}) = \sum_{n \le y} \frac{\Lambda(n)}{n^s \log n} + o(1)$$

if  $s \ge 1/2 + \varepsilon$ . The o(1) terms come from proper prime powers. Since  $\beta = 1 - \xi(u)/\log y \approx 1 - \log u/\log y$ , we certainly have  $s \ge 1/2 + \varepsilon$  if  $y \ge (\log x)^{2+\varepsilon}$ .

In summary: we want to show

$$\sum_{n \leq y} \frac{\Lambda(n)}{n^{\beta} \log n} - I((1-\beta) \log y)$$

<sup>&</sup>lt;sup>1</sup>For simplicity we shall assume  $\sigma \in (1/2, \Theta)$  (instead of  $\sigma \in (1 - \Theta, \Theta)$ ), and concentrate on  $y \geq (\log x)^{2+\varepsilon}$ .

can be 'very' negative if RH fails. Strategy: we fix  $\beta \in (1/2, 1)$ , namely require  $1 - \xi(u) / \log y = \beta$ , which is easy to solve:

$$\xi(u) = \log y(1-\beta) \implies$$
$$e^{\xi(u)} = 1 + u\xi(u) = y^{1-\beta}$$

and

$$1 + u\xi(u) = 1 + u\log y(1 - \beta)$$

 $\mathbf{so}$ 

$$1 + \log x(1 - \beta) = y^{1 - \beta}$$

i.e.

$$y = (1 + \log(1 - \beta))^{1/(1 - \beta)}.$$

Given a function A(x) on  $x \ge 1$ , its Mellin transform is

$$\mathcal{M}A(s) := \int_1^\infty A(x) x^{-s} ds.$$

Landau proved the following.

**Theorem 4.1.** Suppose A(x) is a bounded integrable function on every interval [1, X], which is eventually non-negative. Let  $\sigma_c$  be the infimum of  $\sigma$  such that  $\mathcal{M}A(\sigma)$  converges. Then  $\mathcal{M}A(s)$  is analytic in  $\Re(s) > \sigma_c$  but not at  $s = \sigma_c$ .

To illustrate, let us revisit the proof that  $\psi(x) - x < -x^{\Theta-\varepsilon}$  holds infinitely often, where  $\Theta$  is as before. Consider  $A(x) = \sum_{n \le x} \Lambda(n) - x + x^{\Theta-\varepsilon}$ . Let us suppose A(x) is eventually positive. Not hard to show

$$\mathcal{M}A(s) = -\frac{\zeta'(s-1)}{(s-1)\zeta(s-1)} - \frac{1}{s-2} + \frac{1}{s-1 - \Theta + \varepsilon}.$$

This function is analytic for real  $s > 1 + \Theta - \varepsilon$ , but is not analytic at  $s = 1 + \Theta - \varepsilon$ . Hence, by Landau,  $\mathcal{M}A(s)$  is analytic in the half-plane  $\Re(s) > 1 + \Theta - \varepsilon$ . But this is false – it is only analytic in  $\Re(s) > 1 + \Theta$  due to zeros with real part  $> \Theta - \varepsilon$  for any  $\varepsilon > 0$ ; contradiction.

Another example: Diamond and Pintz (2009) showed

$$\sum_{n \le x} \frac{\Lambda(n)}{n \log n} - \log \log x - \gamma < -\frac{C}{\sqrt{x} \log x}$$

holds infinitely often for any given C > 0, and same with  $> C/(\sqrt{x} \log x)$ . This shows that  $\sqrt{x}(\prod_{p \le x}(1 - 1/p)^{-1} - e^{\gamma} \log x)$  exhibits arbitrarily large positive and negative values as  $x \to \infty$ . They studied the Mellin transform of the LHS.

An almost identical argument works for showing

$$y \mapsto \sum_{n \le y} \frac{\Lambda(n)}{n^{\beta} \log n} - I((1-\beta)\log y) \le -y^{\Theta - \beta - \varepsilon}$$

holds infinitely often.

We conclude that if RH fails, and  $\Theta > 1/2$  is the supremum of the real parts of zeros of  $\zeta$ , then for any  $\beta \in (1/2, \Theta)$  there are sequences  $x_n, y_n$  with  $y_n = (\log x_n)^{1/(1-\beta)+o(1)}$  such that

$$\Psi(x_n, y_n) < x_n \rho(\log x_n / \log y_n) \exp(-y_n^{\Theta - \beta - \varepsilon}).$$

If RH holds,  $\Theta - \beta = 1/2 - \beta < 0$  so this is useless.

**Remark 4.1.** Under RH we can show that  $\Psi(x, y) \sim x\rho(u)F(\beta, y)$  holds for  $y \ge (\log x)^{3/2+\varepsilon}$  and this range is optimal. A similar result holds for polynomials over finite fields, unconditionally.

## 5 Second oscillation result

Finally, let us turn to Theorem 3.1. We assume  $y \leq (\log x)^{2-\varepsilon}$ , so that  $\beta \leq 1/2 - \varepsilon$  (and also  $\alpha \leq 1/2 - \varepsilon$ : it is known that  $\alpha = \beta + O(1/\log y)$ ).

We have seen

$$\frac{\Psi(x,y)}{x\rho(u)} \ll \sqrt{u} \frac{x^{\alpha}\zeta(\alpha,y)}{x^{\beta}\hat{\rho}(-\xi(u))} \ll \sqrt{u} \frac{x^{\beta}\zeta(\beta,y)}{x^{\beta}\hat{\rho}(-\xi(u))} = \sqrt{u} \frac{\zeta(\beta,y)}{\hat{\rho}(-\xi(u))}$$

This used  $\Psi(x,y) \leq x^{\alpha}\zeta(\alpha,y)$ . We also have  $\Psi(x,y) \gg x^{\alpha}\zeta(\alpha,y)/(\sqrt{u}\log y)$  (Hildebrand and Tenenbaum, 80s) if  $y \geq (\log x)^{1+\varepsilon}$ , so

$$\frac{\Psi(x,y)}{x\rho(u)} \gg \frac{x^{\alpha}\zeta(\alpha,y)}{x^{\beta}\hat{\rho}(-\xi(u))\log y} \ge \frac{x^{\alpha}\zeta(\alpha,y)}{x^{\alpha}\hat{\rho}((1-\alpha)\log y)\log y} = \frac{\zeta(\alpha,y)}{\hat{\rho}((1-\alpha)\log y)\log y}$$

The second inequality is trivial (but new): it uses the fact that  $\beta$  minimizes  $s \mapsto x^s \hat{\rho}((1-s)\log y)$ . Recall

$$F(s,y) = \log \zeta(s,y) - \log \hat{\rho}(\log y(s-1)).$$

We have just shown

$$\frac{\Psi(x,y)}{x\rho(u)} \gg e^{F(\alpha,y)}/\log y$$

Unconditionally, Landau's Theorem shows that, if we fix  $\alpha > 0$ ,

$$y \mapsto \sum_{n \le y} \frac{\Lambda(n)}{n^{\alpha} \log n} - I((1-\alpha) \log y)$$

is non-negative. When  $y \leq (\log x)^{2-\varepsilon}$  we have that  $\log F(\alpha, y)$  is larger than  $\sum_{n \leq y} \frac{\Lambda(n)}{n^{\alpha} \log n}$  by a very large quantity, leading to large values of  $\Psi(x, y)/(x\rho(u))$ . Indeed,

$$\log \zeta(s, y) = \sum_{p \le y} -\log(1 - p^{-s}) = \sum_{n \le y} \frac{\Lambda(n)}{n^s \log n} + \sum_{k \ge 2} \sum_{y^{1/k}$$

The k-sum can easily be shown to tend to infinity when  $s \leq 1/2 - \varepsilon$  (this uses nothing more than the Prime Number Theorem), which is the case when  $s = \alpha$  and  $y \leq (\log x)^{2-\varepsilon}$ .